

Cor 6.4 (PNT in arithmetic progressions)

Let ~~...~~ $a \in (\mathbb{Z}/q\mathbb{Z})^\times$. Then,

$$\#\{p \leq x \mid p \equiv a \pmod{q}\} \sim \frac{1}{\varphi(q)} \#\{p \leq x\} \text{ for } x \rightarrow \infty.$$

Pf Let $g(s) := \sum_{\chi} \left(-\frac{L'(s, \chi)}{L(s, \chi)} \right) \cdot \frac{1}{\chi(a)}$

$$= \sum_{n=p^k} \frac{\sum_{\chi} \chi\left(\frac{n}{a}\right) \log p}{n^s}$$

$$= \sum_{\substack{n=p^k: \\ n \equiv a \pmod{q}}} \frac{\varphi(q) \log p}{n^s}$$

It is hol. in $\{\operatorname{Re}(s) \geq 1\}$ except for a simple pole at $s=1$ with residue 1 (coming from $-\frac{L'(s, \chi_0)}{L(s, \chi_0)}$, which comes from the simple pole of $L(s, \chi_0)$ at $s=1$).

Wiener - Ikehara $\Rightarrow \sum_{\substack{n=p^k \leq x: \\ n \equiv a \pmod{q}}} \varphi(q) \log p \sim X.$

$$\parallel \varphi(q) \sum_{\substack{p \leq x: \\ p \equiv a \pmod{q}}} \log p + O(x^{1/2} (\log x)^2)$$

Proceed as ~~the~~ for the PNT (cf. Problem 3 on Ex 1)



Cor 6.5

~~Let $n = x^2 + y^2$...~~

Let $S = \{x^2 + y^2 \mid x, y \in \mathbb{Z}\}$.

We have $\#\{n \in S \mid n \leq x\} \sim C \cdot \frac{x}{\sqrt{\log x}}$

for some constant $C > 0$.

Pf Algebraic NT tells us that $n \in S$ if and only if

n is divisible by each prime $p \equiv 3 \pmod{4}$ an even number of times.

$$\Rightarrow D(1_S, s) = \prod_{p \not\equiv 3 \pmod{4}} \underbrace{\frac{1}{1 - \frac{1}{p^s}}}_{1 + \frac{1}{p^s} + \frac{1}{p^{2s}} + \dots} \cdot \prod_{p \equiv 3 \pmod{4}} \underbrace{\frac{1}{1 - \frac{1}{p^{2s}}}}_{1 + \frac{1}{p^{2s}} + \frac{1}{p^{4s}} + \dots}$$

Let χ_1 be the nontriv. character mod 4.

$$\chi_1(x) = \begin{cases} 1, & x \equiv 1 \pmod{4}, \\ -1, & x \equiv 3 \pmod{4}. \end{cases}$$

$$\begin{aligned} L(s, \chi_1) &= \prod_{p \equiv 1 \pmod{4}} \frac{1}{(1 - \frac{1}{p^s})^2} \cdot \prod_{p \equiv 3 \pmod{4}} \frac{1}{(1 - \frac{1}{p^s})(1 + \frac{1}{p^s})} \\ &= \prod_{p \equiv 1} \frac{1}{(1 - \frac{1}{p^s})^2} \cdot \prod_{p \equiv 3} \frac{1}{1 - \frac{1}{p^{2s}}} \end{aligned}$$

$$\Rightarrow \frac{D(1, s)^2}{L(s, \chi_0)L(s, \chi_1)} = \frac{1}{\left(1 - \frac{1}{2^s}\right)^2} \cdot \prod_{p \equiv 1} 1 \cdot \prod_{p \equiv 3} \frac{1}{1 - \frac{1}{p^{2s}}},$$

which converges for $\operatorname{Re}(s) > \frac{1}{2}$.

$\Rightarrow D(1, s)^2$ ~~is hol.~~ is hol. in $\{\operatorname{Re}(s) \geq 1\}$ except for a simple pole at $s = 1$.

The result follows from Kato's extension of Wiener-Ikehara.

□

7. ~~Functional~~ Functional Equations

We'll generalise the fct. eq. for $S(s)$ to fct. eq. for $L(s, \chi)$.

First, \bullet Poisson summation with a twist:

Lemma 7.1 Let $c: \mathbb{Z}/q\mathbb{Z} \rightarrow \mathbb{C}$ be any function and

let $f \in \mathcal{S}(\mathbb{R})$ (for example). Then,

$$q \cdot \sum_{x \in \mathbb{Z}} c(x \bmod q) f(x) = \sum_{t \in \mathbb{Z}} \hat{c}(t \bmod q) \hat{f}\left(\frac{t}{q}\right)$$

with $\hat{c}: \mathbb{Z}/q\mathbb{Z} \rightarrow \mathbb{C}$ the discrete Fourier transform

$$\text{given by } \hat{c}(t) = \sum_{x \in \mathbb{Z}/q\mathbb{Z}} c(x) e^{2\pi i x t / q}.$$

~~Ex. $c(x) = 1$ $\forall x$ $\Rightarrow \hat{c}(t) = 1$ $\forall t$~~

$$c(x) = 1 \quad \forall x \Rightarrow \hat{c}(t) = \begin{cases} 1, & t = 0 \bmod q, \\ 0, & \text{otherwise.} \end{cases}$$

$\bullet \rightarrow$ claim = Poisson summation.

Pr By linearity, it suffices to consider $c = \mathbb{1}_{\{a \bmod q\}}$.

$$\bullet \text{ Then, LHS} = q \sum_{x \equiv a \bmod q} f(x) = \frac{q}{q} \sum_{\substack{y \in \mathbb{Z} \\ x = qy + a}} g(y)$$

with $g(y) = f(qy + a)$.

$$\text{Poisson summation: } q \sum_y g(y) = q \sum_{t \in \mathbb{Z}} \hat{g}\left(\frac{t}{q}\right) = \sum_{t \in \mathbb{Z}} e^{2\pi i a t / q} \hat{f}\left(\frac{t}{q}\right)$$

= RHS. □

Prmk $\hat{\epsilon}(x) = q \cdot c(-x)$.

~~Prmk~~

primitive

Lemma 7.2 Let χ be a character mod q , extended to $\mathbb{Z}/q\mathbb{Z}$ by 0 (outside $(\mathbb{Z}/q\mathbb{Z})^\times$). Its discrete Fourier transform $\hat{\chi}$

satisfies $\hat{\chi}(\frac{t}{a}) = \overline{\chi(\frac{a}{t})} \cdot \hat{\chi}(1)$. (So the d.f.t. of χ is essentially its complex conjugate.)
for all $t \in \mathbb{Z}/q\mathbb{Z}$.

Def We write $\tau(\chi) := \hat{\chi}(1) = \sum_{x \in (\mathbb{Z}/q\mathbb{Z})^\times} \chi(x) e^{2\pi i x/q}$.
This is called a Gauss sum.

Q Case 1: $t \in (\mathbb{Z}/q\mathbb{Z})^\times$

$$\begin{aligned} \text{LHS} &= \sum_{x \bmod q} \chi(x) e^{2\pi i x t/q} = \sum_{\substack{y \bmod q \\ x=t=y}} \chi\left(\frac{y}{t}\right) e^{2\pi i y/q} \\ &= \overline{\chi(t)} \cdot \sum_{y \bmod q} \chi(y) e^{2\pi i y \cdot 1/q} = \text{RHS.} \end{aligned}$$

$\frac{\chi(y)}{\chi(t)} = \overline{\chi(t)} \cdot \chi(y)$

Case 2: $t \notin (\mathbb{Z}/q\mathbb{Z})^\times$

Let $d = \gcd(t, q)$, $t' = \frac{t}{d}$, $q' = \frac{q}{d}$.

$$\begin{aligned} \text{LHS} &= \sum_{x \bmod q} \chi(x) e^{2\pi i x t'/q'} \\ &= \sum_{x' \bmod q'} \left(\sum_{\substack{x \bmod q \\ x \equiv x' \bmod q'}} \chi(x) \right) e^{2\pi i x' t'/q'} \\ &= 0 = \text{RHS} \quad \text{according to the following Lemma.} \quad \square \end{aligned}$$

Lemma 7.3 Let χ be a ~~primitive~~ character mod q , let

$q' | q$ ~~primitive~~ character mod q' . Then,

$$\sum_{\substack{x \pmod{q}: \\ x \equiv x' \pmod{q'}}} \chi(x) = 0 \quad \text{for all } x' \in (\mathbb{Z}/q'\mathbb{Z})^\times.$$

pf This is clear if $x' \notin (\mathbb{Z}/q'\mathbb{Z})^\times$.

~~Otherwise, take any $x_0 \in (\mathbb{Z}/q'\mathbb{Z})^\times$ with $x_0 \equiv 1 \pmod{q'}$.~~

Otherwise, take any $x_0 \in (\mathbb{Z}/q'\mathbb{Z})^\times$ with $x_0 \equiv 1 \pmod{q'}$.

Mult. by x_0 permutes the summands.

The claim follows, unless $\chi(x_0) = 1$ for all x_0 as above.

But in that case, $\chi(x_1) = \chi(x_2)$ for any $x_1 \equiv x_2 \pmod{q'}$, which implies that χ is induced by a char. of q' , hence not primitive. \square

Cor 7.4 $|\tau(\chi)| = \sqrt{q}$ for any primitive character χ mod q .

pf $\widehat{\chi}(0) = \tau(\chi) \cdot \overline{\chi}(1) \quad \forall \chi$

$$\begin{aligned} \Rightarrow \widehat{\chi}(x) &= \tau(\chi) \cdot \widehat{\chi}(0) = \tau(\chi) \cdot \overline{\chi}(-x) \\ &\stackrel{||}{=} \tau(\chi) \cdot \overline{\chi(-x)} \\ &= \underbrace{\tau(\chi) \cdot \overline{\tau(\chi)}}_{|\tau(\chi)|^2} \cdot \underbrace{\overline{\chi(-x)}}_{\chi(-x)}. \end{aligned}$$

\square