

Math 229 - Introduction to Analytic Number Theory

(Fabian Gundlach)

0. Introduction

A few things that can be proved using analysis:

Thm 0.1 (Prime Number Theorem)

$$\#\{p \leq X \text{ prime}\} \underset{X \rightarrow \infty}{\sim} \frac{X}{\log X}$$

More precise estimate:

$$\#\{p \leq X \text{ prime}\} \sim \int_2^X \frac{1}{\log t} dt.$$

\leadsto Heuristic: The ~~set~~ ^{set} $\{2, 3, 5, \dots\}$ of prime numbers behaves a little like ~~the~~ a random ~~set~~ of natural numbers containing $n \geq 2$ with probability $\frac{1}{\log n}$.

Notation

$$f(x) \sim_{x \rightarrow \infty} g(x):$$

$$\lim_{x \rightarrow \infty} \frac{f(x)}{g(x)} = 1$$

$$f(x) = o_{x \rightarrow \infty}(g(x)):$$

$$\lim_{x \rightarrow \infty} \frac{f(x)}{g(x)} = 0$$

$$f(x) \ll g(x)$$

$$\exists C > 0: \forall x: |f(x)| \leq C \cdot g(x)$$

$$\text{or: } f(x) = o_k(g(x)):$$

$$\forall k: \exists C_k > 0: \forall x: |f(k, x)| \leq C_k \cdot g(k, x)$$

$$f(x) \ll_k g(k, x):$$

$$f(x) \not\sim g(x):$$

$$f(x) \ll g(x) \text{ and } f(x) \not\gg g(x)$$

$$f(x) = \Omega_{x \rightarrow \infty}(g(x)):$$

~~not~~ not $f(x) = o_{x \rightarrow \infty}(g(x))$

$$\text{or: } \limsup_{x \rightarrow \infty} \frac{|f(x)|}{g(x)} > 0.$$

Theorem 0.2 (Dirichlet's Theorem on ~~primes~~ primes in arithmetic progressions)

$$\#\{p \leq X \text{ prime} : p \equiv a \pmod{k}\} \sim \frac{1}{\varphi(k)} \cdot \#\{p \leq X \text{ prime}\}$$

if a is relatively prime to k

~~where~~, where $\varphi(k) = \#(\mathbb{Z}/k\mathbb{Z})^\times$ is the no.

of residue classes $a \pmod{k}$ that are relatively prime to k (invertible). "All invertible res. cl. are equally likely."

E.g., half the primes are $\equiv 1 \pmod{4}$, half are $\equiv 3 \pmod{4}$.

Theorem 0.3

$$\#\{1 \leq n \leq X : \exists a, b \in \mathbb{Z} : n = a^2 + b^2\} \sim C \cdot \frac{X}{\sqrt{\log X}} \text{ for some } C > 0.$$

Theorems 0.1-0.3 are proved using complex analysis (Dirichlet series).

Thm 0.4 (special case of Waring's problem)
Every positive integer is the sum of at most 19
fourth powers.

This is proved using ~~the circle method~~ ^{the} circle method.

Thm 0.5

$$\#\{1 \leq n \leq X \text{ squarefree}\} \sim \frac{6}{\pi^2} \cdot X$$

This is proved using a sieve.

Thm 0.6 (Zhang + polymath)

There are infinitely many pairs of primes that
differ by ~~exactly 2~~ (twin prime conjecture)
at most 246.

Prerequisites: - complex analysis

- Fourier analysis

- a little bit of number theory

Grade: 70% weekly homework (probably due Wednesdays)
(dropping two lowest scores)

30% take-home final exam

OH this week: Mo, Th 3-4pm in room 233

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1. Initiation

1.1. Divisor sum

Def For any integer $n \geq 1$, let $d(n)$ be the number of positive divisors of n :

$$d(n) = \#\{a | n\} = \sum_{a|n} 1.$$

Ex

n	1	2	3	4	5	6
$d(n)$	1	2	2	3	2	4

Goal Estimate $\sum_{n \leq X} d(n)$ for large X .

heuristic

$$\sum_{n \leq X} d(n) = \sum_{n \leq X} \sum_{a|n} 1 = \sum_{\substack{a|b \geq 1: \\ ab \leq X}} 1$$

\uparrow
 $n = ab$

~~$$\sum_{n \leq X} d(n) \approx \sum_{1 \leq a \leq X} \frac{X}{a}$$~~

\uparrow
 $\#\{b: ab \leq X\}$

$$\sum_{n \leq X} d(n) \approx \int_1^X \frac{x}{t} dt = [X \log t]_{t=1}^X = X \log X.$$

Making (I) rigorous:

$$\#\{1 \leq b \leq \frac{x}{a}\} = \lfloor \frac{x}{a} \rfloor = \frac{x}{a} + O(1),$$

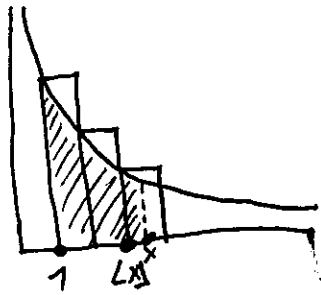
$$\text{so } \sum_{\substack{a|b \geq 1: \\ ab \leq X}} 1 = \sum_{1 \leq a \leq X} \left(\frac{x}{a} + O(1) \right) = \sum_{1 \leq a \leq X} \frac{x}{a} + O(x).$$

Making (I) rigorous:

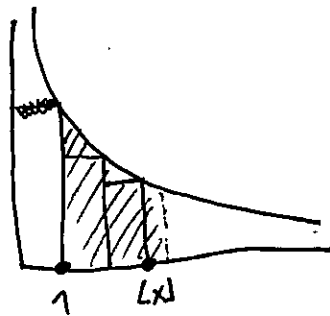
~~claim~~

claim $\sum_{1 \leq n \leq X} \frac{1}{n} = \log X + O(1)$ for $X \geq 1$.

Pf



$$\sum_{1 \leq n \leq X} \frac{1}{n} \geq \int_1^X \frac{1}{t} dt = \log X$$



$$\sum_{2 \leq n \leq X} \frac{1}{n} \leq \int_1^X \frac{1}{t} dt = \log X$$

□

~~□~~

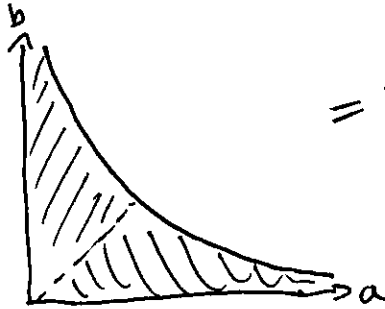
Summary $\sum_{1 \leq n \leq X} d(n) = X \log X + O(X)$,

so the average number of divisors of a random $n \leq X$ is $\sim \log X$ for $X \rightarrow \infty$.

We can improve the estimate!

Improving (I): ("Dirichlet hyperbola method")

$$\sum_{\substack{a, b \geq 1: \\ ab \leq X}} 1 = \sum_{\substack{a \geq b \geq 1: \\ ab \leq X}} 1 + \sum_{\substack{b \geq a \geq 1: \\ ab \leq X}} 1 - \sum_{\substack{a = b \geq 1: \\ ab \leq X}} 1$$



$$= 2 \cdot \sum_{\substack{b \geq a \geq 1: \\ ab \leq X}} 1 - \sum_{\substack{a \geq 1: \\ a^2 \leq X}} 1$$

$$= 2 \cdot \sum_{1 \leq a \leq X^{1/2}} \sum_{a \leq b \leq \frac{X}{a}} 1 - \sum_{1 \leq a \leq X^{1/2}} 1$$

$$= 2 \cdot \sum_{1 \leq a \leq X^{1/2}} \left(\frac{X}{a} - a + O(1) \right) - \left(X^{1/2} + O(1) \right)$$

$$= 2 \cdot \sum_{1 \leq a \leq X^{1/2}} \frac{X}{a} - X + O(X^{1/2})$$

↑
better than $O(X)$

1.2. Abel summation

~~Reminder~~

Reminder (Integration by parts)

Let $f, g: [a, b] \rightarrow \mathbb{C}$ be continuously differentiable. ($a \leq b$)

Then,

$$\int_a^b f'(t)g(t)dt + \int_a^b f(t)g'(t)dt = [f(t)g(t)]_{t=a}^b$$

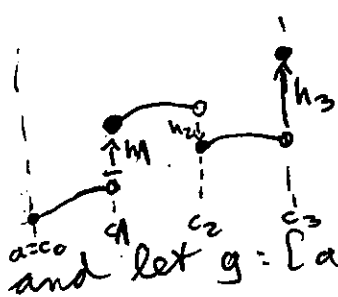
($f(b)g(b) - f(a)g(a)$).

Prblz This continues to hold if f, g are continuous and piecewise continuously differentiable, ignoring points t where $f'(t)$ or $g'(t)$ doesn't exist. It fails if f, g are not continuous.

Thm 1.2.1 (Abel summation)

Let $a = c_0 \leq c_1 \leq \dots \leq c_k = b$, let $f: [a, b] \rightarrow \mathbb{C}$ be continuously differentiable on $[c_i, c_{i+1})$

with a jump of height $h_i = f(c_i) - \lim_{t \nearrow c_i} f(t)$ at c_i



= " $f(c_i) - f(c_i^-)$ " ($i \geq 1$)

↑
limit from below

and let $g: [a, b] \rightarrow \mathbb{C}$ be continuously differentiable.

Then,

$$\int_a^b f'(t)g(t)dt + \sum_{1 \leq i \leq k} h_i g(c_i) + \int_a^b f(t)g'(t)dt = [f(t)g(t)]_{t=a}^b$$

ignore pts, $t = c_i$