

15.2. Dixon's random squares method

Reference: - Chapter 19.5 in
 "Modern Computer Algebra"
 - Dixon: asymptotically fast factorization of integers

Reminder.

Lemma 15.2.1 Let $n = p_1^{e_1} \dots p_k^{e_k}$ odd. Let a be a uniformly random element of $(\mathbb{Z}/n\mathbb{Z})^\times$. Then,

$$1 < \gcd(a-1, n) < n \text{ with probability } 1 - \frac{1}{2^{k-1}} \geq \frac{1}{2}.$$

PP There are exactly 2 bad a :

$$\begin{aligned} a = 1 & : \gcd = n \\ a = -1 & : \gcd = 1. \end{aligned}$$

□

How to construct a \mathbb{Z} -torsion element:

Find $a, b \in (\mathbb{Z}/n\mathbb{Z})^\times$ with $a^2 \equiv b^2 \pmod{n}$. Then, $\frac{a}{b} \in (\mathbb{Z}/n\mathbb{Z})^\times$.

Prime Birthday paradox: need to choose $\sim \sqrt{(\mathbb{Z}/n\mathbb{Z})^\times}$ random elements $a_i \in (\mathbb{Z}/n\mathbb{Z})^\times$ until finding two with the same square.

Idea Try to find numbers $a_1, \dots, a_r \pmod{n}$ s.t.

$$\underbrace{(a_1^2 \pmod{n}) \dots (a_r^2 \pmod{n})}_{\in \{1, \dots, n\}} = b^2 \text{ for an integer } b.$$

This is equivalent to the condition that the LHS is divisible by every prime q an even number of times.

We'll only allow a_i such that $(a_i^2 \pmod{n}) \in \{1, \dots, n\}$ has only prime factors $q < B$.

(small)

alg 15.2.2 ~~the~~ $\frac{\log n}{\log 2} + (\log \log n)(25+2)$

Given: odd integer $n \geq 3$ ^{composite,} ~~not a prime power,~~

some $B \geq 2$. $\frac{\log n}{\log B} \approx \log_B n$

1) compute the primes $p_1, \dots, p_r \leq B$.

Lemma 15.2.2 ~~Let~~ Let $n = p_1^{e_1} \dots p_k^{e_k}$ odd, $k \geq 2$.

Let $q_1 < \dots < q_r$ be ^(small) primes numbers ~~not dividing~~ [not dividing n].

The following alg. returns ~~a uniformly random element of~~ ~~satisfies~~ ~~with~~ ~~probability~~ ~~at least~~ ~~1/2~~

a uniformly random element of $(\mathbb{Z}/n\mathbb{Z})^{\times [2]}$.

alg (15.2.2)

Find uniformly random elements $b_1, \dots, b_{r+1} \in (\mathbb{Z}/n\mathbb{Z})^{\times}$ ^{subjs} ~~to the condition~~ ~~that~~ $(b_i^2 \pmod n)$ ~~each~~ ~~is~~ ~~coprime~~ ~~to~~ ~~n~~

~~can~~ can be written as $(b_i^2 \pmod n) = q_1^{f_{i1}} \dots q_r^{f_{ir}}$:

Just pick random $b \in (\mathbb{Z}/n\mathbb{Z})^{\times}$, compute $b^2 \pmod n$, find the number of times this integer is divisible by ~~each~~ each q_i by trial division, until you've found $r+1$ ~~good~~ ~~good~~ residue classes b_i . [This terminates because $b=1$ is possible.]

2) Using Gaussian elimination, find ~~the~~ the kernel of the map $\mathbb{F}_2^{r+1} \rightarrow \mathbb{F}_2^r$

$(v_i)_{i=1, \dots, r+1} \mapsto (\sum_i v_i f_{ij})_{j=1, \dots, r}$

3) Pick a uniformly random ~~element~~ nonzero element $v = (v_i) \in \mathbb{F}_2^{r+1}$ of the kernel. Let $S = \{i \mid v_i = 1\} \subseteq \{1, \dots, r+1\}$.

$$\left(\Rightarrow \sum_{i \in S} v_i f_{ij} \equiv 0 \pmod{2} \text{ for all } j. \right)$$

4) Let $t_j = \frac{1}{2} \sum_{i \in S} f_{ij}$.

$$\left(\Rightarrow \prod_{i \in S} (b_i^2 \pmod{n}) = \prod_{i \in S} \left(\prod_j q_j^{f_{ij}} \right) = \prod_j q_j^{2t_j} = \left(\prod_j q_j^{t_j} \right)^2 \right)$$

5) Return $c = \frac{\prod_{i \in S} b_i}{\prod_j q_j^{t_j}} \pmod{n}$.

Pf The result c is ~~an el.~~ an el. of $(\mathbb{Z}/n\mathbb{Z})^\times(\mathbb{Z})$ because

$$\left(\prod_{i \in S} b_i \right)^2 \equiv \left(\prod_j q_j^{t_j} \right)^2 \pmod{n}.$$

~~We'll show that for any fixed set $S \neq \emptyset$, all elements of $(\mathbb{Z}/n\mathbb{Z})^\times(\mathbb{Z})$ are equally likely. Fix any $i_0 \in S$.~~

~~Fix the set S and any $i_0 \in S$ and the value~~

~~$d = (b_{i_0}^2 \pmod{n})$. Also fix all b_j with $j \neq i_0$.~~

~~$\Rightarrow b_{i_0}$ is determined by its square d up to mult. by an el. of $(\mathbb{Z}/n\mathbb{Z})^\times$. All square roots of $d \pmod{n}$ are equally likely to be the value of b_{i_0} .~~

~~\Rightarrow all elements of $(\mathbb{Z}/n\mathbb{Z})^\times(\mathbb{Z})$ occur with the same probability.~~



~~We'll show that for any fixed S, c_0, c_1, c_2~~

We'll show that c is a uniformly random element of $(\mathbb{Z}/n\mathbb{Z})^{\times} [2]$, even ~~for any fixed~~

~~$S, \{a_i \in S, b_i \text{ for all } i \neq i_0, d_{i_0} = (b_{i_0}^2 \text{ mod } n)$~~

~~Note that S can be determined from~~

particular fixed values $d_i = (b_i^2 \text{ mod } n)$.

Note that, S only depends on these values (and randomness)
 the set

not on the square roots b_i of d_i .

• The b_i are uniformly random square roots of the d_i .

→ ~~the~~ $\prod_{i \in S} b_i$ is a uniformly distributed random square root of $\prod_{i \in S} d_i$ (even if we pick $i_0 \in S$ and fix all b_i with $i \neq i_0$).

Principle We could have chosen v (and therefore S) deterministically, as long as the choice only depends on d_1, \dots, d_{r+1} , not on b_1, \dots, b_{r+1} .

Question ~~What fraction~~ ^{For} what fraction of elements $b \in (\mathbb{Z}/n\mathbb{Z})^\times$ can $(b^2 \pmod n)$ be written as $q_1^{f_1} \dots q_r^{f_r}$?
 i.e. how long does step 1 take?

Proof Assume $q_1 < \dots < q_r$ and $q_r^t \leq n$. Then,

$$\# \{1 \leq a \leq n \mid a = q_1^{f_1} \dots q_r^{f_r} \text{ for some } f_1, \dots, f_r \geq 0\}$$

$$\geq \# \{(f_1, \dots, f_r) \mid f_1, \dots, f_r \geq 0, f_1 + \dots + f_r \leq t\}$$

$$= \binom{t+r}{t} \geq \frac{r^t}{t!} \quad (\text{"close to } \frac{n}{t!})$$

But we need to ~~show~~ prove that many of these $1 \leq a \leq n$ are quadratic residues mod n .

invertible Idea ~~show~~ (quadr. nonres. mod p_i) · (quadr. nonres) = (quadr. res)

Lemma 15.2.3 ~~Assume~~ ~~Assume~~ ~~Assume~~ $q_1 < \dots < q_r$ and $q_r^{2t} \leq n$

and that no q_i divides n . Then,

~~$$\# \{a \in (\mathbb{Z}/n\mathbb{Z})^\times \mid \dots\}$$~~

~~(quadr. res)~~

~~$$\# \{a \in \mathbb{Z}/n\mathbb{Z} \mid a = q_1^{f_1} \dots q_r^{f_r} \text{ for some } f_1, \dots, f_r \geq 0\}$$~~

$$\# \{b \in (\mathbb{Z}/n\mathbb{Z})^\times \mid (b^2 \pmod n) = q_1^{f_1} \dots q_r^{f_r} \text{ for some } f_1, \dots, f_r \geq 0\}$$

$$\geq \frac{r^{2t}}{(2t)!}$$

("close to $\frac{n}{(2t)!}$ ")

write $n = p_1^{e_1} \dots p_r^{e_r}$. $\mathbb{Z}/n\mathbb{Z} \cong C_{\varphi(p_1^{e_1})} \times \dots \times C_{\varphi(p_r^{e_r})}$.

pf consider the map

$$\chi: (\mathbb{Z}/n\mathbb{Z})^\times \rightarrow (\mathbb{Z}/n\mathbb{Z})^\times / (\mathbb{Z}/n\mathbb{Z})^{\times 2} \cong \underbrace{C_2 \times \dots \times C_2}_k =: G$$

↑
quadr. res

with kernel $(\mathbb{Z}/n\mathbb{Z})^{\times 2}$.

For any $g \in G$, let

$$U_g := \chi^{-1}(g).$$

If $a_1, a_2 \in U_g$, then $\chi(a_1) = \chi(a_2) = g$, so $a_1 a_2 \in \ker = (\mathbb{Z}/n\mathbb{Z})^{\times 2}$ is a quadratic residue, with 2^k square roots.

Let $T := \{a \in (\mathbb{Z}/n\mathbb{Z})^\times \mid a = q_1^{f_1} \dots q_r^{f_r} \text{ for some } f_1, \dots, f_r \geq 0\}$
with $f_1 + \dots + f_r \leq 2t$
($\Rightarrow 1 \leq a \leq \sqrt{n}$)

If $a_1, a_2 \in U_g \cap T$, then

$$a_1 a_2 \in W := \left\{ 1 \leq a \leq n \mid \begin{array}{l} a = q_1^{f_1} \dots q_r^{f_r} \\ f_1 + \dots + f_r \leq 2t \\ a \in (\mathbb{Z}/n\mathbb{Z})^{\times 2} \end{array} \right\}.$$

Hence, we obtain a map

$$\rho: \bigsqcup_{g \in G} (U_g \cap T) \times (U_g \cap T) \longrightarrow W.$$

$(a_1, a_2) \longmapsto a_1 a_2$

[crude estimate:]

any $a = q_1^{f_1} \dots q_r^{f_r} \in W$ (with $f_1 + \dots + f_r \leq 2t$)

has at most $\binom{2t}{t} = \frac{(2t)!}{t!^2}$ preimages (choose which of the $2t$ prime factors of a go into a_1).

\Rightarrow ~~scribble~~

$$\sum_{g \in G} |U_{g \cap T}|^2 \leq |W| \cdot \frac{(2t)!}{t!^2}$$

~~scribble~~

AM-QM ineq: $\left(\frac{\sum_{g \in G} |U_{g \cap T}|^2}{|G|} \right)^2 \leq \frac{\sum_{g \in G} |U_{g \cap T}|^2}{|G|}$

~~scribble~~ $(2/n)^x = \prod_{g \in G} U_g \rightarrow \left\| \frac{|T|}{|G|} \right\|^2$

$|W| \cdot \frac{(2t)!}{t!^2}$

~~scribble~~

\Rightarrow ~~scribble~~ $\# \{b \in (2/n)^x \mid (b^2 \text{ mod } n) = q_1^{f_1} \dots q_r^{f_r} \dots\}$

$$\geq 2^k \cdot |W| \geq \frac{2^k \cdot |T|^2}{|G| \cdot (2t)! / t!^2} = \frac{2^k}{\binom{t+r-1}{t}} \cdot \frac{t!^2}{(2t)!}$$

$$\geq \frac{2^k}{t!^2} \cdot \frac{t!^2}{(2t)!} = \frac{2^k}{(2t)!}$$

□