

F.2. Rank

an $n \times n$ -matrix

Rule The rank of M is the largest $0 \leq r \leq n$ s.t. some $r \times r$ -minor of M (made from r not necessarily consecutive rows and columns) has nonzero determinant.

Cor F.2.1 $\text{rk}(M) \geq \text{rk}(M \bmod p)$ by primop

with equality if p doesn't divide ~~any $r \times r$ -minor~~ the (nonzero) det of a ~~particular~~ $r \times r$ -minor of M , where $r = \text{rk}(M)$.

Cor F.2.2

$$\text{rk}(M) = \max_{p \in N} \text{rk}(M \bmod p) \text{ if } \prod_{p \in N} p \mid M$$

$$> T(M) := \prod_{p \in N} \left(\frac{\text{rk}(M \bmod p)}{\left(\sum_{j=1}^n m_{ij}^2 \right)^{1/2}} \right) \quad (\text{Blaugemins})$$

which can be computed in time

$$O((n + \log T(M)) \cdot n^w)$$

Rule If $\prod_{p \in N} p \mid M$ and $N' \geq N$, then the probability

that a random prime $p \leq N'$ doesn't satisfy *

$$\text{rk}(M) = \text{rk}(M \bmod p)$$

is at most $\frac{\#\{p \leq N\}}{\#\{p \leq N'\}}$.

This gives rise to a Monte Carlo alg. with

$$\text{running time } O(n^w + \underbrace{(n + \log T(M)) \cdot \log \log(n + \log T(M)))}_{\text{time to find primes } p \leq n + \log T(M)})$$

time to find primes $p \leq n + \log T(M)$

7.3. Resultants

Rule: Let $f, g \in \mathbb{Q}(x)$ be ~~more pol.~~ ~~relatively prime in $\mathbb{Q}(x)$~~

~~so that~~ If $f, g \bmod p$ are relatively prime in $\mathbb{F}_p[x]$,
then f, g are rel. prime in $\mathbb{Q}(x)$.

The converse doesn't hold:

~~x^2+1 is irreduc. in $\mathbb{Q}[x]$, but $x^2+1 = (x+1)^2 \bmod 2$.~~

~~Could there be something like~~

E.g. $x^2+1, x+1$ are rel. prime in $\mathbb{Q}(x)$, but $x^2+1 = (x+1)^2 \bmod 2$

Q If f, g are rel. prime over \mathbb{Q} , for which p are they
not rel. prime ~~mod p~~?

~~Def~~ For any $d > 0$, let $K[X]_d := \{f \in K[x] : \deg(f) < d\}$.

Lemma 7.3.1 Let $f, g \in K[x]$ be ~~pol.~~ of degrees n, m .

Then, ~~gcd(f, g) = 1 if and only if~~

the map $\{ \text{fact}(f) : \deg(f) = m \} \times \{ \text{fact}(g) : \deg(g) = n \} \rightarrow \{ \text{fact}(fg) : \deg(fg) = n+m \}$

$$(a, b) \mapsto f_a + g_b$$

is an isomorphism.

Q Note that $\dim(LHS) = m+n = \dim(RHS)$.

$$K[x]_{cm} \cong K[x]/f, \quad K[x]_{cn} \cong K[x]/g, \quad K[x]_{c_{n+m}} \cong K[x]/fg.$$

" \Leftarrow " ~~If $\gcd(f, g) = e$ then fg is not constant, then the image~~
only contains multiples of $\gcd(f, g)$. \Rightarrow The map isn't
surjective.

" \Rightarrow " The map is an isom. according to ~~Bézout's identity~~ Bézout's identity. \square

Def The resultant $\text{Res}^{(f,g)}$ of pol. $f, g \in K[x]$ of deg. n, m is the determinant of the map in Lemma 7.3.1 w.r.t the basis $((1, 0), (x, 0), \dots, (x^{m-1}, 0), (0, 1), \dots, (0, x^{n-1}))$ of the LHS & the basis $((1, x, \dots, x^{n+m-1}))$ of the RHS.

Cor 7.3.2 $\gcd(f, g) = 1 \Leftrightarrow \text{Res}(f, g) \neq 0$.

Cor 7.3.3 Let $f, g \in Q(x)$ be pol. and let p be a prime not dividing the denominator of any coeff. of f or g . Then, f, g are rel. prime mod p iff $p \nmid \text{Res}(f, g)$.

orange box

$$f(x) = x^n + a_{n-1}x^{n-1} + \dots + a_0 \text{ and}$$

$$g(x) = x^m + b_{m-1}x^{m-1} + \dots + b_0$$

Show $\text{Res}(f, g) = \det$

$$\begin{matrix} a_0 & 0 & b_0 & 0 \\ \vdots & \ddots & \vdots & \ddots & b_0 \\ a_{n-1} & a_0 & b_{m-1} & b_0 \\ 0 & \vdots & 0 & \vdots \\ a_{n-m} & 0 & b_{m-n} & b_0 \end{matrix} \quad n+m$$

Sylvester matrix

Lemma 7.3.03

a) $\text{Res}(f, g) = (-1)^{nm} \text{Res}(f, g)$

c) $\text{Res}(f, g) = a_n^m b_m^n \cdot \prod_{\substack{1 \leq i \leq n \\ 1 \leq j \leq m}} (\alpha_i - \beta_j) = a_n^m \prod_{i=1}^n g(\alpha_i)$

if $\alpha_1, \dots, \alpha_n \in \bar{K}$ are the roots of f (with mult.)

and $\beta_1, \dots, \beta_m \in \bar{K}$

b) $\text{Res}(rf, sg) = r^m s^n \text{Res}(f, g) \quad \forall r, s \in K^\times$

of a), b) clear

c) w.l.o.g. f and g are monic: $a_n = b_m = 1$.

$$\Rightarrow f(x) = \prod_i (x - \alpha_i), \quad g(x) = \prod_j (x - \beta_j)$$

\Rightarrow l.h.s. of $\text{Res}(f, g)$ is hom. pol. in $\alpha_1, \dots, \alpha_n$ of deg. $n-k$.

-- b_k of g -- -- β_1, \dots, β_m -- -- $m-l$.

$\Rightarrow \text{Res}(f, g)$ is hom. pol. in $\alpha_1, \dots, \alpha_n, \beta_1, \dots, \beta_m$ of deg. nm .

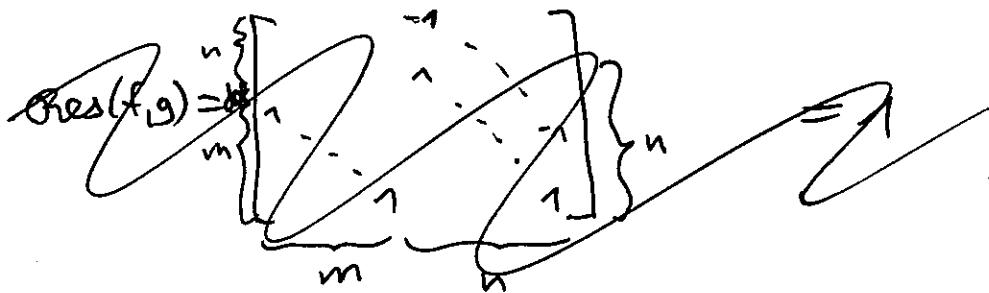
expand the determinant

$$\begin{aligned} & \text{If } (\alpha_i)_{i \in I} \in \bar{K}^n, (\beta_j)_{j \in J} \in \bar{K}^m \\ & \text{satisfy } \alpha_i = \beta_j \text{ for some } i, j, \text{ then } X - \alpha_i \mid f, g, \text{ so } \gcd(f, g) \neq 1, \text{ so } \text{Res}(f, g) = 0. \\ & \Rightarrow \underbrace{\prod_{i,j} (\alpha_i - \beta_j)}_{\deg nm} \text{ divides } \underbrace{\text{Res}(f, g)}_{\deg nm} \text{ as a pol. in } \alpha_1, \dots, \alpha_n, \beta_1, \dots, \beta_m \end{aligned}$$

$$\Rightarrow \text{Res}(f, g) = C_n \cdot \prod_{i,j} (\alpha_i - \beta_j) \text{ for some constant } C_n.$$

To show $C_{n,m} = 1$, it suffices to check the equality for one pair (f, g) of pol. f, g of deg. n, m .

For example, look at $f(x) = x^n$, $g(x) = x^m + 1$



$$\alpha_1 = \dots = \alpha_n = 0, \beta_j$$

$$\text{Res}(f, g) = \det \begin{bmatrix} \alpha_1 & \dots & \alpha_n \\ \vdots & \ddots & \vdots \\ 1 & \dots & 1 \end{bmatrix} = \cancel{\dots} / 1$$

$$\alpha_1 = \dots = \alpha_n = 0$$

$$\Rightarrow \prod_{i,j} (\alpha_i - \beta_j) = \underbrace{\left(\prod_j (-\beta_j) \right)^n}_{\text{const. coeff. of } g} = \cancel{\dots} 1.$$

□

Point Resultants can be computed using the ^(fast) Euclidean algorithm. (HW)

over fields K
with $O(1)$
arithmetic

- * The CRT trick then allows us to compute resultants of polynomials in $\mathbb{Q}(X)$, in part. to determine whether two pol. in $\mathbb{Q}[X]$ are relatively prime.