

We will focus on imaginary quadratic number fields K .
 (Zahner, McLurley: § rigorous subexponential alg. for computation of class group
 (For ~~the~~ general number fields K , see
~~the~~ Buchmann: § subexp. alg. for the determination of
 class groups and regulators of abs. nr. fields).

Prmk \bullet $r_1 = 0$, $r_2 = 1$, $\mathcal{O}_K^\times = \mu_K$,
 $R_K = 1$,
 $w_K = |\mu_K| = \begin{cases} 2, & \bullet \text{ otherwise} \\ 4, & K = \mathbb{Q}(i) \\ 6, & K = \mathbb{Q}(\zeta_3) \end{cases}$

$K = \mathbb{Q}(\sqrt{D_K})$, \bullet $D_K < 0$.

~~Def~~

~~An ideal $0 \neq \mathfrak{a} \subset \mathcal{O}_K$ is reduced if ~~the~~~~

~~a) it is not divisible by any integer $n \geq 2$ and~~

~~b)~~

(Nonstandard)

Def A fractional ideal \mathfrak{a} is reduced if ~~there is no~~
~~nonzero element of minimal norm~~
 $1 \in \mathfrak{a}$ but there is no $x \in \mathfrak{a}$ with $Nm(x) < 1$.
" | compl. emb. (x)

Prmk $1 \in \mathfrak{a} \Leftrightarrow \mathcal{O}_K \subseteq \mathfrak{a} \Leftrightarrow \mathfrak{a}^{-1} \subseteq \mathcal{O}_K$,
" (1)

so ~~the~~ the inverse of a reduced fractional ideal is an (integral) ideal of \mathcal{O}_K .

Lemma 16.3.3 any ideal class contains ^{at least 1, at most 6} ~~at least 1, at most 6~~ reduced ideals

Pf Let \mathfrak{b} be any fractional ideal ~~and let~~
~~and let~~ and let $y \in \mathfrak{b}$ be a nonzero element of minimal norm.
Then, $\mathfrak{a} = y^{-1} \cdot \mathfrak{b}$ is reduced. any reduced ideal in the ideal class $[\mathfrak{b}]$ is
~~of this form~~ of this form. There are at most 6 such y . □

Prmk ~~we can efficiently determine the reduced ideals in the same ideal class as~~
~~we can efficiently determine the reduced ideals in the same ideal class as~~
We can efficiently determine the reduced ideals using Gauss's lattice reduction to find the shortest, nonzero vector. Hence, we can efficiently determine whether two ideals lie in the same ideal class.

Thm 16.3.4 (Minkowski bound)
If \mathfrak{a} is reduced, then $Nm(\mathfrak{a}^{-1}) \leq \mathcal{O}(\sqrt{|D_K|})$.

Prmk This gives rise to a slow alg. to determine \mathcal{O}_K :
Find all ideals \mathfrak{b} with $Nm(\mathfrak{b}) \leq \mathcal{O}(\sqrt{|D_K|})$.
For each, compute the ~~reduced~~ reduced ideals in the same ideal class as \mathfrak{b}^{-1} to determine which \mathfrak{b} are in the same class.

Thm 16.3.5 Assume the ~~extended~~ ^{extended} Riemann hypothesis ~~(ERH)~~ ^(ERH).
 Then, ll_K is generated by the (ideal classes of) prime ideals \mathfrak{p}
~~of norm~~ $N_{\mathbb{Q}}(\mathfrak{p}) \leq 6 (\log |D_K|)^2$.

Proof Hence, if $\mathfrak{p}_1, \dots, \mathfrak{p}_r$ are the prime ideals of norm $\leq B$
 with $B \geq 6 (\log |D_K|)^2$, then we get a surjective group hom

$$\varphi: \mathbb{Z}^r \longrightarrow \text{ll}_K$$

$$(a_1, \dots, a_r) \mapsto [\mathfrak{p}_1^{a_1} \cdots \mathfrak{p}_r^{a_r}]$$

To determine ll_K , we need to find its kernel

$$\text{ll}_K \cong \mathbb{Z}^r / \ker(\varphi)$$

~~the~~ i.e.: Need to find elements generating the rank n lattice $\ker(\varphi)$.

Idea: Find random elements ~~of $\ker(\varphi)$~~ until they generate $\ker(\varphi)$.

How to tell when we're finished?

Let $\Lambda \subset \ker(\varphi)$, the lattice generated by the elements discovered so far.

~~We either have $\Lambda = \ker(\varphi)$,~~

~~or~~

$\exists \Lambda \subsetneq \ker(\varphi)$, then $[\ker(\varphi) : \Lambda] \geq 2$, so

$$|\mathbb{Z}^r / \Lambda| \geq 2 \cdot |\mathbb{Z}^r / \ker(\varphi)| = 2 |\text{ll}_K|.$$

Hence, it suffices to know $|\text{ll}_K|$ within a factor of 2, which (assuming ERH) can be ~~computed~~ computed using the class number formula.

How to find ^{"quickly"} random elements of $\ker(\varphi)$?

Pick a random vector $a = (a_1, \dots, a_r) \in \mathbb{Z}^r$.

~~Compute~~

The ideal $\langle \varphi_1^{a_1}, \dots, \varphi_r^{a_r} \rangle$ is

[This only lies in $\ker(\varphi)$ with prob. $\approx \frac{1}{\#\text{ll}_u}$ ^(very small)!]

Compute ^{fractional} a reduced ideal α in the ~~same~~ ideal class $[\varphi_1^{a_1}, \dots, \varphi_r^{a_r}] = [\varphi_1]^{a_1} \dots [\varphi_r]^{a_r}$ using fast exponentiation, reducing at every step to ~~ensure~~ ensure that we only need to work with ideals of norm $\leq |D_u|$ at any time.

If $\alpha^{-1} = \varphi_1^{b_1} \dots \varphi_r^{b_r}$ with integers $b_1, \dots, b_r \geq 0$,

then \bullet ~~compute~~ $[\varphi_1^{a_1}, \dots, \varphi_r^{a_r}, \varphi_1^{b_1}, \dots, \varphi_r^{b_r}] = [\alpha \alpha^{-1}] = [1]$,

so \bullet $[a+b] = (a_1 + b_1, \dots, a_r + b_r) \in \ker(\varphi)$.

(Note that $b_1, \dots, b_r \leq O(\log |D_u|)$

because $\text{Nm}(\alpha^{-1}) \leq O(\sqrt{|D_u|})$.)

Otherwise, try ~~another~~ random vector $a \in \mathbb{Z}^r$.
again, with a new

~~First to obtain linearly independent~~

~~pick the first r vectors~~

(uniformly)

We pick the first r vectors a^V from

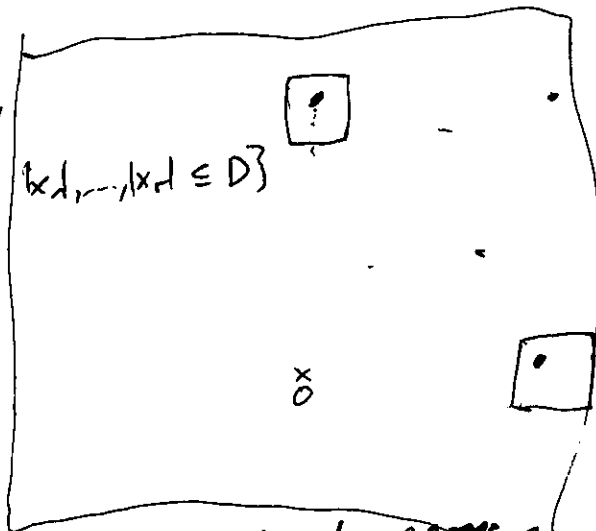
$$(2r|D|, 0, \dots, 0) + \mathcal{B}(D),$$

$$\{x \in \mathbb{Z}^r \mid |x_1|, \dots, |x_r| \leq D\}$$

$$(0, 2r|D|, \dots, 0) + \mathcal{B}(D),$$

\vdots

$$(0, \dots, 0, 2r|D|) + \mathcal{B}(D)$$



so the ~~the~~ first r elements of $\text{ker}(\varphi)$ we construct ~~space~~ a random

lattice $\Lambda \subset \mathbb{Z}^r$ of covolume $\leq \mathcal{O}((r|D|)^r)$.

ranked and

afterwards, we pick vectors a from $\mathcal{B}(|D|^2)$ until $|\mathbb{Z}^r/\Lambda|$ is small enough.

Analyzing the ~~expected~~ running time and choosing \mathcal{B} optimally, we get:

bound on norm. \mathcal{B} : we consider

Thm 16.3.6 (~~real~~ ~~McLure~~)

Assuming GRH, we can determine l_n

in expected time $\mathcal{O}(\exp(\sqrt{r} \cdot \sqrt{\log |D| \log \log |D|^r}))$.

↑
can presumably be improved...